

Ex. 1.

1) $M \models \forall x \exists y (f(x,y) = a)$

iff for each member $b \in Q, M \models \neg (b/x)$

iff for each member $b \in Q,$

there's at least one element $b_1 \in Q, M \models \neg_1(b, b_1)$

iff for each member $b \in Q, there's $b_1 \in Q, such that $(+(b, b_1), 1) \in <, or in other notation.$$$

$b + b_1 < 1.$

it holds by the additive inverse existence of rational numbers.

$\varphi = \forall x \exists y (f(x,y) = a)$

$= \forall x \neg (x)$

$= \forall x \exists y (\neg_1(x,y))$

$= \forall x \exists y (= (f(x,y), a))$

$= \forall x \exists y (\neg_2(x,y))$

2) $M \models \forall x \neg R(x,a)$

iff for $b \in Q, M \models \neg (b/x)$

iff for $b \in Q, M \not\models \neg_1(b/x)$

iff for $b \in Q, (b, 1) \notin <, or in another notation $b \not< 1$$

$\varphi = \forall x \neg (x)$

$= \forall x \neg_1(x)$

But $1/2 \in Q$ and $1/2 < 1$. it follows $M \not\models \forall x \neg R(x,a)$

3) $M \models \forall x (R(b/a) \rightarrow (R(f(x,b), f(x,a))))$

iff for $b \in Q, M \models \neg (b/x)$

iff for $b \in Q, Either $M \not\models \neg_1(b/x)$ or $M \models \neg_2(b/x)$$

$\varphi = \forall x \neg (x)$

$= \forall x \neg_1(x) \rightarrow \neg_2(x)$

but $M \not\models \neg_1(b/x)$ is equivalent to $2 \not< 1$ which holds.

Then φ is satisfied in M

4.) Call $\varphi(z) = \forall x \forall y (f(x, y) = z)$

$$\begin{aligned}\varphi(1/2) &= \forall x \forall y (f(x, y) = 1/2) \\ &= \forall x \neg \exists y (f(x, y) = 1/2) \\ &= \forall x \forall y \neg (f(x, y) = 1/2)\end{aligned}$$

We show it's not satisfiable on any member $r \in \mathbb{Q}$,
Concluding $\mathcal{M} \not\models \varphi$

$$\mathcal{M} \models \varphi(r/z)$$

iff for $b \in \mathbb{Q}$, $\mathcal{M} \models \neg (b/x, r/z)$

iff for $b, b_1 \in \mathbb{Q}$, $\mathcal{M} \models \neg (b/x, b_1/y, r/z)$

iff for $b, b_1 \in \mathbb{Q}$, $=^{\mathcal{M}} (f^{\mathcal{M}}(b/x, b_1/y), z^{\mathcal{M}})$, or in another notation $b + b_1 = r$

But it doesn't hold. For example $r/4 \in \mathbb{Q}$ and $r/4 + r/4 \neq r$
Hence Not satisfiable.

Ex. 2

Note. The convention for the first-order languages is to contain equality symbol "="

$$1) \varphi = \forall x (R(0, x) \vee = (0, x))$$

$$\begin{aligned} \varphi^N &= \forall x (R^N(0, x) \vee =^N(0, x)) \\ &= \forall x (x > 0 \vee 0 = x) \end{aligned}$$

$$2) \varphi(x, y) = \exists k = (g(x, k), y)$$

$$\begin{aligned} \varphi^N(x, y) &= \exists k =^N (g^N(x, k), y) \\ &= \exists k \cdot x \cdot k = y \end{aligned}$$

$$3) \psi(x, y, z) = \left(\left(\varphi(z, x) \wedge \varphi(z, y) \right) \wedge \left(\forall k (\varphi(k, x) \wedge \varphi(k, y)) \rightarrow (= (k, z) \vee R(k, z)) \right) \right)$$

$$\psi^N(x, y, z)$$

$$\begin{aligned} &= \left(\left(\varphi^N(z, x) \wedge \varphi^N(z, y) \right) \wedge \left(\forall k (\varphi^N(k, x) \wedge \varphi^N(k, y)) \rightarrow (=^N(k, z) \vee R^N(k, z)) \right) \right) \\ &= \left((\exists k \cdot z \cdot k = x \wedge \exists k \cdot z \cdot k = y) \wedge (\forall k (\exists z \cdot k \cdot z = x \wedge \exists z \cdot k \cdot z = y) \rightarrow k = z \vee k < z) \right) \end{aligned}$$

$$4) \theta(x) = \left(\forall k \varphi(k, x) \rightarrow (= (k, x) \vee = (k, b)) \right)$$

$$\begin{aligned} \theta^N(x) &= \left(\forall k \varphi^N(k, x) \rightarrow (=^N(k, x) \vee =^N(k, b^N)) \right) \\ &= \left(\forall k \exists z \cdot k \cdot z = x \rightarrow (k = x \vee k = 1) \right) \end{aligned}$$

Ex. 3

Vector addition associativity

$$\forall u, \forall v, \forall w = \left(g(u, g(v, w)), g(g(u, v), w) \right)$$

$$u + (v + w) = (u + v) + w$$

Vector addition commutativity

$$\forall u, \forall v = (g(u, v), g(v, u))$$

$$u + v = v + u$$

Vector addition identity

$$\exists v, \forall u, g(f_0^V(v), u) = g(u, f_0^V(v)) = u$$

$$\exists v \in V, \forall u \in V, 0 + u = u + 0 = u$$

$$\text{alt. } \exists r \in R, \exists v, \forall u, g(f_r^V(v), u) = g(u, f_r^V(v)) = u$$

Vector addition inverse

$$\forall v, \exists u, g(v, u) = f_0^V(u)$$

$$\forall v \in V, \exists u \in V, v + u = 0$$

Scalar multiplication compatibility

$$\text{for any } a \text{ and } b \text{ in } R, \forall v, f_a^V(f_b^V(v)) = f_{ab}^V(v)$$

$$\forall v \in V, a, b \in F, a(b \cdot v) = (ab) \cdot v$$

Scalar multiplication identity

$$\forall v, f_1^V(v) = v$$

$$\forall v \in V, 1 \cdot v = v$$

Scalar multiplication distributivity

$$\text{for any } a \text{ in } R, \forall v, \forall u, f_a^V(g(u, v)) = g(f_a^V(u), f_a^V(v))$$

$$\forall u, v \in V, \forall a \in F, a(u + v) = au + av$$

$$\text{for any } a \text{ and } b \text{ in } R, \forall v, f_{a+b}^V(v) = g(f_a^V(v), f_b^V(v))$$

$$\forall v \in V, \forall a, b \in F, (a + b)v = av + bv$$

The interpretation of these formulas under a structure is satisfied if and only if the structure satisfies real vector spaces axioms.

For example, Consider $\mathcal{V} = (V, 0, \oplus, \otimes)$, Satisfying for any $v \in V$.

$$\ast 0 \oplus v = v \oplus 0 = v \quad (1)$$

$$\ast 0 \otimes v = 0 \quad (2)$$

For \mathcal{L} -formula $\Psi = \exists v \forall u g(f_0^V(v), u) = g(u, f_0^V(v))$, Observe

$$\mathcal{V} \models \Psi$$

iff

there's a member $b \in V$, such that for all $b_1 \in V$,

$$0 \otimes b + b_1 = b_1 + 0 \otimes b$$

$$0 + b_1 = b_1 + 0 = b_1$$

Ex. 4

We disprove by a counter-example.

Consider the structure $\mathcal{N} = (N, I)$ where N is the set of natural numbers and I is a unary identity function $I: x \mapsto x$.

Let as hinted $\varphi(x) = \exists y \neg (x=y)$.

Clearly $\mathcal{N} \models \forall x \varphi(x)$, i.e. for every natural number there's a distinct natural number from it.

Consider the term $T = g(y)$, with function symbol g and variable symbol y , and the formula $\varphi(T) = \exists y \neg (T=y) = \exists y \neg (g(y)=y)$.

Whereby $g^{\mathcal{N}}$ is the identity I , $\varphi(T)$ is not satisfied under structure \mathcal{N} . By let the interpretation means there's a natural number b such that $I(y) \neq y$, Contradiction.

In conclusion, $\mathcal{N} \not\models \varphi(T)$ for the term T .

We are given for any member $b \in \Gamma$, $\Gamma \models \varphi(b/x)$.

Consider an arbitrary term T .

Case ① T is a constant

Then T^M is a member of Γ .

Hence $M \models \varphi(T^M/x)$.

Case ② T is a variable x_i

by hypothesis we know $M \models (b/x)$ for any $b \in \Gamma$

← Plain
Possibly Variables

Case ③ $T = f(t_1, \dots, t_n)$ for some terms t_1, \dots, t_n

Then $T^M = f^M(t_1^M, \dots, t_n^M) \in \Gamma$, A member of Γ

Hence $M \models \varphi(T^M/x)$ \Leftarrow Not well-defined

From Case ①, Case ②, and Case ③, $M \models \varphi(T)$ for any term T .