

Ex. 1

lemma. For any formulas  $\alpha, \beta, \gamma$ .

- 1.  $\alpha \vdash (\beta \rightarrow \alpha)$
- 2.  $\{(\gamma \rightarrow \alpha), (\alpha \rightarrow \beta)\} \vdash (\gamma \rightarrow \beta)$
- 3.  $\{(\alpha \rightarrow (\beta \rightarrow \gamma)), \beta\} \vdash (\alpha \rightarrow \gamma)$
- 4.  $\vdash (x = x)$
- 5.  $\vdash (x = y \rightarrow y = x)$

proof.

- 1.  $\alpha$  Assumption  
 $(\alpha \rightarrow (\beta \rightarrow \alpha))$  Ax. 1  
 $(\beta \rightarrow \alpha)$  MP
  
- 2.  $((\gamma \rightarrow (\alpha \rightarrow \beta)) \rightarrow ((\gamma \rightarrow \alpha) \rightarrow (\gamma \rightarrow \beta)))$  Ax. 2  
 $(\alpha \rightarrow \beta)$  Assumption  
 $(\gamma \rightarrow (\alpha \rightarrow \beta))$  lemma 1  
 $((\gamma \rightarrow \alpha) \rightarrow (\gamma \rightarrow \beta))$  MP  
 $(\gamma \rightarrow \alpha)$  Assumption  
 $(\gamma \rightarrow \beta)$  MP

$$3. \{ \beta, \alpha \rightarrow (\beta \rightarrow \gamma) \} \vdash (\alpha \rightarrow \gamma)$$

$$(\alpha \rightarrow (\beta \rightarrow \gamma)) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow \gamma)) \quad \text{Ax. 2}$$

 $\beta$ 

Assumption

$$(\alpha \rightarrow \beta)$$

lemma 1

$$(\alpha \rightarrow (\beta \rightarrow \gamma))$$

Assumption

$$((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow \gamma))$$

MP

$$(\alpha \rightarrow \gamma)$$

MP

4.

Given  $\vdash (t=t)$  in p. 84, it immediately follows  $\vdash (x=x)$ .

$$5. \vdash (x=y \rightarrow y=x)$$

$$x=y \rightarrow ((x=x) \rightarrow (y=x))$$

Ax. 7

$$(x=x)$$

lemma 4

$$(x=y \rightarrow y=x)$$

lemma 3

1.)

$$(y=x \rightarrow (y=z \rightarrow x=z))$$

Ax. 7

$$(x=y \rightarrow y=x)$$

lemma 5

$$(x=y \rightarrow (y=z \rightarrow x=z))$$

lemma 2

2)

Claim.  $\{ \forall x \forall y \varphi \} \vdash \forall y \forall x \varphi$

$\forall x \forall y \varphi$

Assump.

$\forall x \forall y \varphi \rightarrow \forall y \varphi(x)$

Ax. 4

$\forall y \varphi(x)$

MP

$\forall y \varphi(x) \rightarrow \varphi(x, y)$

Ax. 4

$\varphi(x, y)$

MP

$\forall x \varphi(y)$

Gen

$\forall y \forall x \varphi$

Gen

Since all occurrences of  $x$  and  $y$  in  $\forall x \forall y \varphi$  are bound, we can use the deduction theorem to conclude

$\vdash (\forall x \forall y \varphi \rightarrow \forall y \forall x \varphi)$  ■

Ex. 2

$M \models (\forall x(\varphi \rightarrow \tau) \rightarrow (\varphi \rightarrow \forall x\tau))$ , where  $x$  isn't free in  $\varphi$

iff if  $M \models \forall x(\varphi \rightarrow \tau)$ , and if  $M \models \varphi$ , Then  $M \models \forall x\tau$

iff if

\* for any  $b \in M$ , if  $M \models \varphi$ , then  $M \models \tau(b/x)$ . Observe we substituted  $\varphi(x)$  by  $\varphi$  as  $x$  is not free in  $\varphi$ . They're the same formulas.

\*  $M \models \varphi$ .

Then  $M \models \tau(b/x)$  for any  $b \in M$ .

The last statement clearly holds. ■

Ex. 3.

This is exactly the contrapositive of thm 4.4.5.

Ex. 4.

1.) We prove it by induction along the pairing axiom.

Claim. For all  $n \geq 1$ , given  $X_1, X_2, \dots, X_n$  are sets, so is  $\bigcup_{i=1}^n X_i$ .

Base Case.  $n=1$ .  $\bigcup_{i=1}^1 X_i = X_1$ , already given as a set.

Induction step. For  $n > 1$ , we know both  $\bigcup_{i=1}^{n-1} X_i$  and  $X_n$  are sets.

Then  $\bigcup_{i=1}^n X_i = \bigcup_{i=1}^{n-1} X_i \cup X_n$  is a set by pairing axiom.

2.) Already given in lemma 5.1.1.